



DIALOG

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# Namespace workshop

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# One typical problem

Workspace 1

Run  
DATA

Workspace 2

Go  
DATA





# One typical problem

Workspace

Run  
DATA

Go  
DATA





# One typical problem

Workspace

Run  
WS1\_DATA

Go  
WS2\_DATA





# One typical problem

Workspace

WS1\_Run  
WS1\_DATA

WS2\_Go  
WS2\_DATA





# One typical problem

Workspace

WS1 namespace

WS1.Run  
WS1.DATA

WS2 namespace

WS2.Go  
WS2.DATA





# One typical problem

Workspace

WS1 namespace

Run  
DATA

WS2 namespace

Go  
DATA





# One typical problem

Workspace 1

Run  
DATA

Workspace 2

Go  
DATA





# Using namespaces

Tree	Workspace	File system
Node	Namespace	Directory
Leaf	Name	File
Roots	# □SE	/ or C: D: E:
Separator	.	/ or \
Parent Node	##	..
Current Node	□THIS	.
Create Node	□NS	mkdir
Change Node	□CS	cd
Create Leaf	←	>

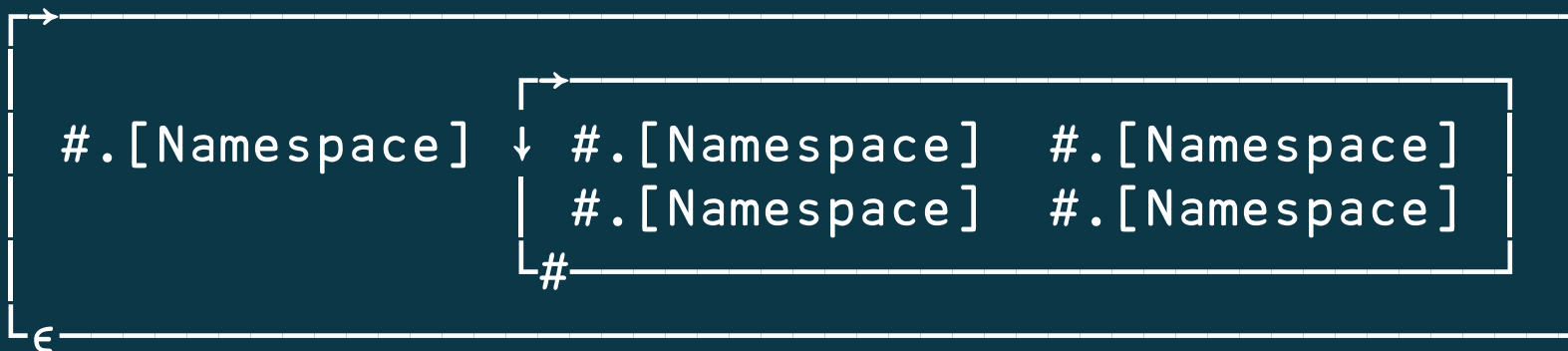




# Using namespaces

`ns←⎕NS ''`      a create a namespace

`]display ns (2 2⍴ns)`      a it's a new kind of scalar





# Using namespaces

```
data←7 8 9
ns←⎕NS ' '
ns.data←'hello'
⎕←(data)(ns.data)
7 8 9 hello
⎕CS ns      ⌘ change current space
⎕←data
hello
⎕CS #       ⌘ come back to root
⎕←data
7 8 9
```





# Using namespaces

Exercise (easy)

Write a function that tells you whether a namespace is a root

$$\{\omega \in \# \mid \square SE\}$$
$$\{\omega \equiv \omega.##\}$$




# Using namespaces

Exercise (medium)

Write a function that returns the root of a namespace

$$\{p \leftarrow \omega.## \quad \diamond \quad \omega \equiv p:\omega \quad \diamond \quad \nabla p\}$$
$$\{\omega.##\}^* \equiv$$




# Dotted expressions

Left side must be an array of namespace(s)

```
(ns1 ns2)←(⊞NS'')(⊞NS'')
(ns1.num ns2.num)←77 99
⊞←(ns1 ns2).num
77 99
⊞←(2 3pns1 ns2).num
77 99 77
99 77 99
(ns1 ns2).num←55 ♦ ⊞←(ns1 ns2).num
55 55                a scalar extension
```





# Dotted expressions

Right side may be any APL expression (without  $\diamond$ )  
executed in the namespace(s) on the left side

```
(ns1 ns2)←(⊞NS'')(⊞NS'')
ns1.(a b)←'hello' 'world' ⋄ ⊞←ns1.(a,b)
helloworld
ns2.a←{2/ω} ⋄ ns2.b←7 8 9 ⋄ ⊞←ns2.(a,b)
7 7 8 8 9 9
⊞←(2 3pns1 ns2).(a,b) ♪ which languages allow this?
helloworld 7 7 8 8 9 9 helloworld
7 7 8 8 9 9 helloworld 7 7 8 8 9 9
```





# Dotted expressions

Exercise (easy)

What's the difference between  $NS . A \leftarrow B$  and  $(NS . A) \leftarrow B$  ?

None

What's the difference between  $NS . A \leftarrow B$  and  $NS . (A \leftarrow B)$  ?

B is looked up in NS (even if it is local!)

The whole expression  $(A \leftarrow B)$  is evaluated in NS





# Dotted expressions

Exercise (easy)

Write a function that returns the roots of an arbitrary array of namespaces

The previous answer still works

$\{\omega . \#\# \} \dot{*} \equiv$





# By-value VS by-reference

Arrays are interpreted by value

```

n ← 7 8 9 ♦ m ← n ♦ □ ← n m
7 8 9 7 8 9      n arrays have the same value

```

```

n[2] ← 80 ♦ □ ← n m
7 80 9 7 8 9      n m still has the same value

```

```

m[3] ← 90 ♦ □ ← n m
7 80 9 7 8 90     n still has the same value

```





# By-value VS by-reference

Arrays are interpreted by value

```
n ← 3 ♦ f ← n ◦ + ♦ □ ← f 100  
103
```

```
n ← 4 ♦ □ ← f 100  
103      ∩ n was passed by value - f hasn't changed
```

```
f ← n ◦ + ♦ □ ← f 100  
104      ∩ f has changed - not n
```





# By-value VS by-reference

Arrays are interpreted by value

```
▽ foo arg
  arg[2]←10
▽
n←7 8 9 ♦ foo n ♦ □←n
7 8 9   n n hasn't changed
▽ arg←foo arg
  arg[2]←10
▽
n←foo n
```





# By-value VS by-reference

Exercise (medium)

Is there a difference between  $(\Box NS ' ')(\Box NS ' ')$  and  $(2\rho\Box NS ' ')$ ?

By-value: no, there is only one empty namespace:  $\theta\theta \equiv 2\rho\Box NS$

By-reference: yes, in the first case we create two entities,  
in the second case only one

When should two namespaces compare equal?

By-value: when they happen to have the exact same content (slow)

By-reference: when they originate from the same call to  $\Box NS$  (fast)





# By-value VS by-reference

Namespaces are interpreted by reference

A namespace is an identifiable container, irrespectively of the content

An array is a conceptual value, irrespectively of the actual instance

In the context of Dyalog APL, namespaces ARE called “references” or “refs”

They have a different name classes (2=array ; 3=function ; 9=reference)





# By-value VS by-reference

Namespaces are interpreted by reference

```

ns1←[]NS ''
ns2←ns1      ρ both names designate the same namespace
[]←ns2.vec
VALUE ERROR  ρ name undefined yet
ns1.vec←7 8 9 ♦ []←ns2.vec
7 8 9        ρ name defined
ns2.vec[2]←10 ♦ []←ns1.vec
7 10 9       ρ value has changed

```





# By-value VS by-reference

Namespaces are interpreted by reference

```

▽ new goo arg
  arg.vec[2]←new
▽
(ns1←[]NS'').vec←7 8 9 ♦ 10 goo ns1 ♦ []←ns1.vec
7 10 9                               ρ namespace has been modified
ns2←ns1 ♦ 100 goo ns1 ♦ []←ns2.vec
7 100 9                              ρ namespace has been modified
1000 goo ns3←[]NS'' ♦ []←ns3.vec
VALUE ERROR                          ρ vec is not defined in ns3

```





# By-value VS by-reference

Namespaces are interpreted by reference

```

▽ foo arg
  arg[2]←10
▽
(ns1←[]NS'').vec←7 8 9
foo ns1.vec ♦ []←ns1.vec
7 8 9          ρ ns1.vec is an array passed by value

```





# By-value VS by-reference

Namespaces are interpreted by reference

```
(ns1←[]NS '').vec←7 8 9
ns3←[]NS ns1      ρ take a deep copy
ns3.vec[2]←10  ♦  []←(ns1 ns3).vec
7 8 9  7 10 9      ρ only the new copy has changed
```

Namespaces by value are rare but possible 😊





# By-value VS by-reference

Namespaces semantics are by reference, and not by value (assignment, comparison, argument passing...) unless you really want to.

This allows

- tracking individual entities independently of their content
- pass modifiable arguments to functions (use with care)

Because of the different semantics, it is worth distinguishing names of references. Depth-0 references are of name class 9, depth $\geq 1$  arrays of refs are of name class 2

Recognise dottable names





# By-value VS by-reference

Exercise (easy)

Write a function that copies a workspace into a namespace

$\{\alpha. \square \text{CY } \omega\}$

NB. No need to return a result - caller already knows  $\alpha$





# By-value VS by-reference

Exercise (medium)

Make it work on an arbitrary array of namespaces

Requires care to distribute the one workspace name to multiple namespaces

$\{\alpha. \Box CY \text{ } c \ddot{*} (\equiv \omega) \vdash \omega\}$        $\mu$  for positive depth

$\{0 \equiv \alpha : \alpha. \Box CY \text{ } \omega \diamond \alpha \nabla \ddot{''} c \omega\}$        $\mu$  for negative depth

$\{\alpha. (\Box CY \text{ } \omega)\}$        $\mu$  Brenner's trick





# By-value VS by-reference

Exercise (medium)

Write a boolean-returning function that detects namespaces

Tip: The name class of a scalar namespace is 9. The name class of non-scalar arrays and of non-namespace scalars is 2.

```
IsScalarRef ← { 9 = NC ' ω ' }
RefMask ← { 0 ≡ ω : 9 = NC ' ω ' } ♦ ∇ `` ω }
```

Homework : RefMask could use  $\square$ DR for performance





# Parent hierarchy (optional)

Namespaces have a single immutable parent, fixed at creation time

```
ns←[]NS'' ♦ []←#=ns.##  
1          ρ created in #  
ns.sub1←ns.[]NS'' ♦ []←ns=ns.sub1.##  
1          ρ created in #.ns  
ns.sub1.(sub2←[]NS'') ♦ []←ns.sub1=ns.sub1.sub2.##  
1          ρ created in #.ns.sub1  
[]←ns.sub1.sub2.##.##.##  
#
```





# Parent hierarchy (optional)

Names of a namespace are NOT (necessarily) its children  
Children are NOT (necessarily) named from their parent

```

ns0←[]NS'' ♦ ns1←[]NS'' ♦ ns1.ns2←ns0
[]←ns1=ns1.ns2.##
0 ρ ns1.ns2 is ns0 which was created in #
  ns3←ns1.ns2.[]NS'' ♦ []←ns0=ns3.##
1 ρ yet ns0 contains no name designating its child ns3

```

Just like arrays, namespaces do not know their names  
(they only know their parent)





# Parent hierarchy (optional)

Exercise (hard) :

Write a function that lists the children of a namespace

Tip 1. `ΩNL` gives the list of names, not the list of children

Tip 2. impossible without crawling through the whole workspace





# Parent hierarchy (optional)

```

▽ children←{arg}ListChildren target;args;next;parents;visited
  :If 0=□NC'arg' ♦ arg←(# □SE target)(0ρ#)(0ρ#) ♦ :EndIf
  (parents children visited)←arg
  next←uεparents.(⌵``'##' '□THIS',□NL ⁂9)⌵ visit all reachable ns
  next~←visited ⌵ excepted visited ones
  :If 0εpnext ♦ :Return ♦ :EndIf ⌵ (0ρ#).## is NONCE ERROR
  childrenu←(target=next.##)/next ⌵ append children of target
  visited,←next ⌵ no next has been visited
  :If 0εpnext ♦ :Return ♦ :EndIf ⌵ F``0ρ# is NONCE ERROR
  args←next children visited
  childrenu←args ListChildren target ⌵ recur on unvisited ns

```

▽





# Display form (optional)

Since namespaces may or may not have a name,  
their display form is ARBITRARY.

```
ns←⎕NS ' ' ⋄ ⎕←2ρns
#. [Namespace]   #. [Namespace]
  ns.⎕DF '<My Namespace>' ⋄ ⎕←2ρns
<My Namespace>   <My Namespace>
```

Just like arrays, namespaces do not know their names.





# Display form (optional)

Crafting namespaces where  $(\phi \circ \tau)$  is identity requires care.  
All parents must be correctly named.

```
'named'⌵NS''
'named.sub'⌵NS''  ρ better than named.('sub'⌵NS'')
ns1←named ⋄ ns2←named.sub ⋄ ⌵←ns1 ns2
#.named   #.named.sub
⌵←{ω≡ϕτω}⌵ns1 ns2
1 1
```

Namespaces are then a proper tree where each node knows its path.





# Display form (optional)

Exercise (medium)

Write a function that lists the children of a namespace, assuming they're all named

$$\{\omega.(\Phi^{\cdot}\Box NL^{-9})\}$$

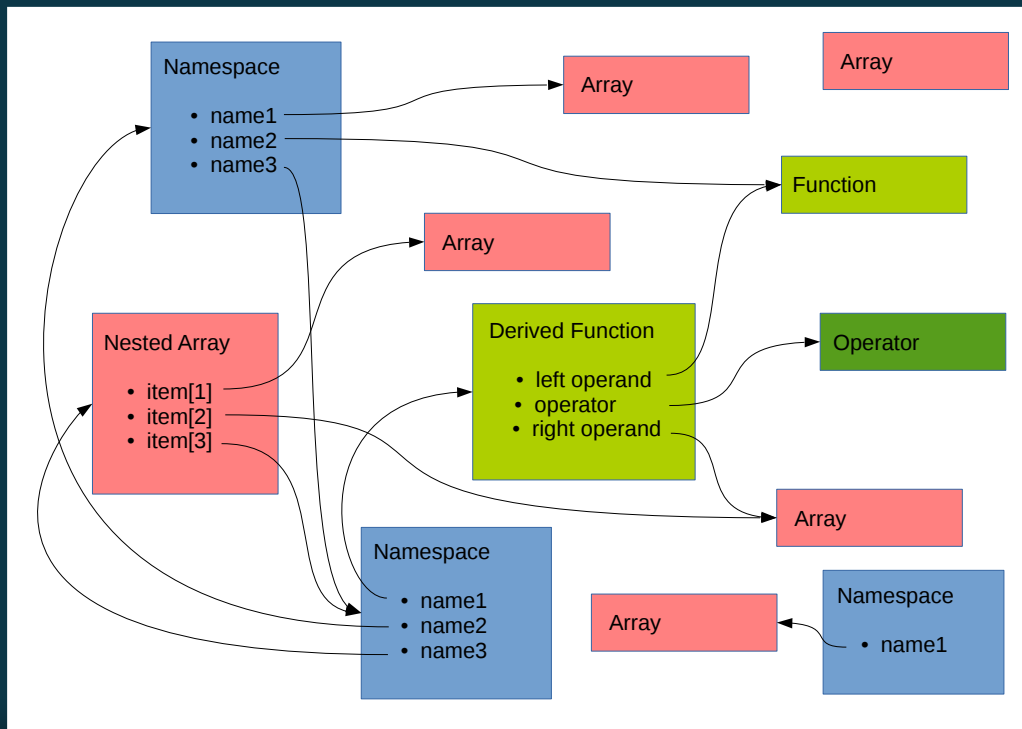
Homework: supplementary filtering with `STATE` might be required if the namespace may have a function on the stack with locals of name class 9

Hack of the day: one can turn a name class 9 (scalar reference) into a name class 2 by ravelling it





# The workspace fauna





# The workspace fauna

A namespace contains nothing but

- a list of names:  $\omega . \square \text{NL} - \imath 10$
- a reference to its parent:  $\omega . \# \#$
- a local copy of system variables:  $\square \text{IO} \ \square \text{CT} \ \square \text{RL} \dots$
- a display form:  $\omega . \square \text{DF} \rightarrow \rightarrow \ \mathfrak{f} \omega$

A name can designate any entity (array, function, operator, namespace)

Just like arrays, namespaces do NOT know their names.

Namespaces do NOT contain entities

Names of a namespace are NOT children of the namespace





# Two approaches to namespaces

Usage	Creation	Display Form	Parent Hierarchy
Grounded	Named	Full Path	Tree
Floating	Unnamed	Description	Flat

Each approach is easy, mixing them is tough !





# Scripted namespaces

Store a namespace as a single piece of text

Helpful for text-based version control

```

)ed *NS      a equivalent to '*ED'NS'
[]FIX  ':Namespace NS'  'VAR←123'  ':EndNamespace'
[]←[]SRC NS
  
```

Every time the script is fixed, then the namespace is cleared and the script is re-run. This happens only once at run-time, but many times at development time.





# At the doorstep of Object Orientation

Objects ARE namespaces, with a few constraints:

- Can only call functions and modify variables
- Cannot create variables or change code
- New tree hierarchy called “derived classes” to avoid duplicating code
- Possibility to “hide” internal code

This is the fashionable way to provide an API







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