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Static Analysis of APL for Tooling and Compliance

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Static Analysis

***** *****

When running your code isn't enough

“Reading your code”

Checking

Changing

Tool Support for Static Analysis

Getting you started now...

And a peek at the future

Make it easier to answer
questions about your code

What kind of questions, I hear you say?

Questions about your code

- Do I have singletons living in my code that might surprise me later?
- Do I have any code that depends on certain errors? (Complex, 128-bit floats, and returning nested results)
- Does my code handle empty array prototypes (Reduce, Each, Rank) correctly?
- Have I sanitized my data coming from unverified sources?
- Have I ensured data security between my client data?
- Do I have hidden dependencies on index origin?
- Do I index all my database arrays using appropriate field names?

More questions...

- Am I using some old idioms that can be replaced by newer primitives?
- Do I have naked strand assignments?
- Are my strand assignments vulnerable to being confused for modified assignments?
- Do I make dangerous assumptions about empty arrays that aren't justified?
- Can I convert functions back and forth between trad-fns and dfns safely?
- Am I passing an operator a function that unintentionally captures shadowed names in the operator?
- Do I have any strand assignments against vectors that might not be the right length?

Even more questions...

- Do I have any rank or type errors sitting around?
- Do have code whose empty prototype has the wrong element type? E.g., " vs θ
- Is this loop iteration independent (I can parallelize it)?
- Is it safe to reorder these statements?
- Does this function set any of my locally shadowed bindings?
- Do I have any duplicate names or labels?
- Where am I using a function via alias instead of via namespace reference? $F \leftarrow ns.F$

Even even more questions....

- Do I have any race conditions between threads?
- Do I depend on variables that might not exist?
- Do I have dead or useless code that is sitting around unused or secretly broken?
- Are there unused variables in my function that I have shadowed?
- Do I have unshadowed variables that really should be shadowed?
- Do I bind a local variable in a dfn but use it before it is shadowed?

Crystal clear, eh?

Let's go looking for trouble...

Looking for trouble...

```
+/(QUERY_MASK≠DB)[;1]
```

```
+/(QUERY_MASK≠DB)[;1]  
      ↑  
      1
```

1. Magic numbers for column indexing

`+/(QUERY_MASK≠DB)[;AGES]`

↑

1

1. Magic numbers for column indexing

`+/(QUERY_MASK/DB)[;AGES]`

 ↑ ↑

 2 1

1. Magic numbers for column indexing
2. What if we want to convert to an inverted table?

`+ / ( QUERY_MASK / DB ) [ ; AGES ]`

 ↑ ↑ ↑

 3 2 1

1. Magic numbers for column indexing
2. What if we want to convert to an inverted table?
3. Is this a filter or a reduction?

$+ \neq (\text{QUERY_MASK} \neq \text{DB}) [; \text{AGES}]$

 ↑ ↑ ↑ ↑

 4 3 2 1

1. Magic numbers for column indexing
2. What if we want to convert to an inverted table?
3. Is this a filter or a reduction?
4. What if $\text{QUERY_MASK} \neq$ were a complex expression?

~~+~~ (QUERY_MASK ~~≠~~ DB) [; AGES]

↑ ↑ ↑ ↑ ↑

5 4 3 2 1

1. Magic numbers for column indexing
2. What if we want to convert to an inverted table?
3. Is this a filter or a reduction?
4. What if QUERY_MASK ~~≠~~ were a complex expression?
5. ...

```
DB←3 2p'brandon' 2 'aaron' 12
QUERY_MASK←DB[;0]ε'brandon' 'aaron'
(QUERY_MASK≠DB)[;1]
```

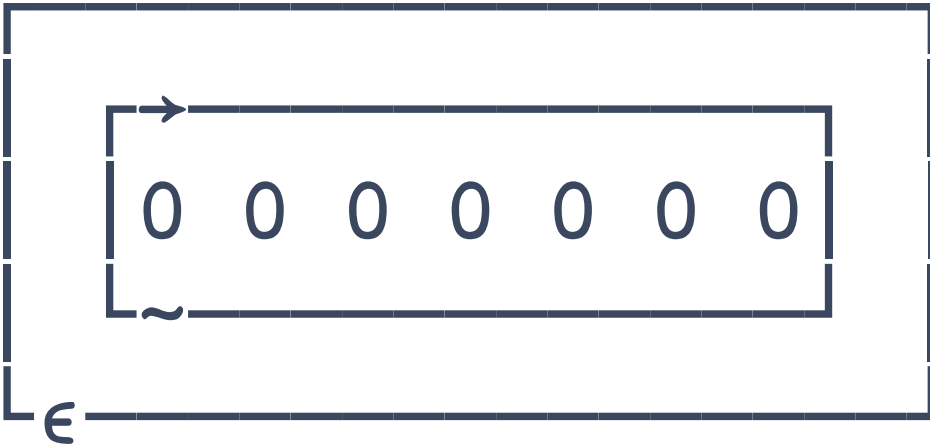
2 12 2

```
+≠(QUERY_MASK≠DB)[;1]
```

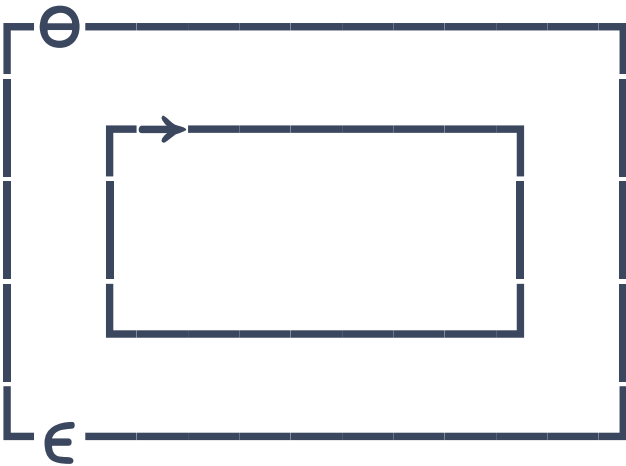
16

```
QUERY_MASK ← DB[;0] ∈ c'morten'  
+ / (QUERY_MASK / DB)[;1]
```

QUERY_MASK ← DB[;0] € c'morten'
+ / (QUERY_MASK / DB)[;1]



QUERY_MASK ← DB[;0] € c'morten'
(QUERY_MASK ≠ DB)[;1]



`+ / ( QUERY_MASK / DB ) [ ; AGES ]`

↑ ↑ ↑ ↑ ↑
5 4 3 2 1

1. Magic numbers for column indexing
2. What if we want to convert to an inverted table?
3. Is this a filter or a reduction?
4. What if `QUERY_MASK /` were a complex expression?
5. **Empty-array prototypes must be considered**

Find this code...

$+ \neq (\text{QUERY_MASK} \neq \text{DATABASE}) [\text{; AGES}]$

```
(ast _ symbol) ← 3 ↑ codfns.PS src  
(parent _ type kind name _ vardef _ _) ← ast  
  
expr ← type ∈ Strand Expr  
roots ←  $\perp$  parent =  $\perp$  ≠ parent  
dexpr ← (type = Expr) ∧ kind = 2  
moxpr ← (type = Oper) ∧ kind ∈ 1 2  
defs ← (type = Var) ∧ type[parent] = Header  
namespace ← (type = Trad) ∧ kind = 0  
ns ← { $\alpha[\omega]$ } @ {~namespace[ $\omega$ ]} *  $\equiv$  ~parent  
prmfn ← (type = Prim) ∧ kind = 2  
prmop ← (type = Prim) ∧ kind = 3  
named ← {name = symbol |  $\iota \subset, \omega$ }
```

```
:Namespace  
  DB←LD'DYNA'  
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]  
  +/(QUERY_MASK≠DB)[;AGES]  
:EndNamespace
```

```
db←_defs^(ns∈roots)^named'DB'
```

```
:Namespace  
  DB←LD'DYNA'  
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]  
  +/(QUERY_MASK≠DB)[;AGES]  
:EndNamespace
```

```
db←ldefs^(ns∈roots)^named'DB'
```

:Namespace

- $DB \leftarrow LD 'DYNA'$
- $QUERY_MASK \leftarrow DB[;EMAIL] \in (LD 'DYALOG')[;EMAIL]$
- $+ \neg (QUERY_MASK \neq DB)[;AGES]$

:EndNamespace

$db \leftarrow \underline{1} \text{ defs } \wedge (ns \in roots) \wedge \text{named} 'DB'$

```
:Namespace  
  DB←LD'DYNA'  
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]  
  +/(QUERY_MASK≠DB)[;AGES]  
:EndNamespace
```

```
db←_defs^(ns∈roots)^named'DB'
```

:Namespace

DB←LD'DYNA'

QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]

+/(QUERY_MASK≠DB)[;AGES]

:EndNamespace

db←ldefs^(ns∈roots)^named'DB'

```
:Namespace  
  DB←LD'DYNA'  
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]  
  +/(QUERY_MASK≠DB)[;AGES]  
:EndNamespace
```

```
db←_defs^(ns∈roots)^named'DB'  
dbex←(type=Var)^(vardef∈db)^expr[parent]
```

```
:Namespace
  DB←LD'DYNA'
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]
  +/(QUERY_MASK≠DB)[;AGES]
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db←_defs^(ns∈roots)^named'DB'
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  +/(QUERY_MASK≠DB)[;AGES]
:EndNamespace

```

```

db←_ids^(ns∈roots)^named'DB'
dbex←(type=Var)^(vardef∈db)^expr[parent]

```

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db←_defs^(nse roots)^named'DB'  
dbex←(type=Var)^(vardef∈db)^expr[parent]
```

```

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```

```

db←_defs^(ns∈roots)^named'DB'
dbex←(type=Var)^(vardef∈db)^expr[parent]
{dbex[parent]v←dbex^expr[parent]}*≡θ

```

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  DB←LD'DYNA'
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db←_defs^(nse roots)^named'DB'
dbex←(type=Var)^(vardef∈db)^expr[parent]
{dbex[parent]v←dbex^expr[parent]}*≡θ
idxpr←parent[_prmf n^(named'[')^dexpr[parent]]

```

```

:Namespace
  DB ← LD 'DYNA'
  QUERY_MASK ← DB [ ; EMAIL ] ∈ ( LD 'DYALOG' ) [ ; EMAIL ]
  + / ( QUERY_MASK / DB ) [ ; AGES ]
:EndNamespace

```

```

db ← ldefs ^ ( ns ∈ roots ) ^ named ' DB '
dbex ← ( type = Var ) ^ ( vardef ∈ db ) ^ expr [ parent ]
{ dbex [ parent ] v ← dbex ^ expr [ parent ] } * ≡ θ
idxpr ← parent [ lprmf n ^ ( named ' [ ' ) ^ dexpr [ parent ] ]

```

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idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]

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idxpr←parent[_prmf n^(named'[')^dexpr[parent]]
idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]

```

```

:Namespace
  DB ← LD 'DYNA'
  QUERY_MASK ← DB[;EMAIL] ∈ (LD 'DYALOG')[;EMAIL]
  + / (QUERY_MASK / DB)[;AGES]
:EndNamespace

```

```

db ← ldefs ^ (ns ∈ roots) ^ named 'DB'
dbex ← (type = Var) ^ (vardef ∈ db) ^ expr[parent]
{dbex[parent] v ← dbex ^ expr[parent]} * ≡ θ
idxpr ← parent[lprmf n ^ (named ' ' ) ^ dexpr[parent]]
idxpr_db ← parent[ldbex ^ (≠parent) ^ parent ∈ idxpr]

```

```

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  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]
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idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]

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:Namespace
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```

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dbex←(type=Var)^(vardef∈db)^expr[parent]
{dbex[parent]v←dbex^expr[parent]}*≡θ
idxpr←parent[_prmf n^(named'[')^dexpr[parent]]
idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]
redox←parent[_prmop^(named'≠')^moxpr[parent]]

```

```

:Namespace
  DB←LD'DYNA'
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]
  +/(QUERY_MASK≠DB)[;AGES]
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idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]
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```

```

:Namespace
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:EndNamespace

```

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idxpr←parent[_prmf n^(named'[')^dexpr[parent]]
idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]
redox←parent[_prmop^(named'≠')^moxpr[parent]]

```

```

:Namespace
  DB←LD'DYNA'
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]
  +/(QUERY_MASK≠DB)[;AGES]
:EndNamespace

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db←_defs^(nse roots)^named'DB'
dbex←(type=Var)^(vardef∈db)^expr[parent]
{dbex[parent]v←dbex^expr[parent]}*≡θ
idxpr←parent[_prmf n^(named'[')^dexpr[parent]]
idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]
redox←parent[_prmop^(named'≠')^moxpr[parent]]
parent[redox]∩parent[idxpr_db]

```

```

:Namespace
  DB←LD'DYNA'
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]
  +/ (QUERY_MASK≠DB)[;AGES]
:EndNamespace

```

```

db←_defs^(nse roots)^named'DB'
dbex←(type=Var)^(vardef∈db)^expr[parent]
{dbex[parent]v←dbex^expr[parent]}*≡θ
idxpr←parent[_prmf n^(named'[')^dexpr[parent]]
idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]
redox←parent[_prmop^(named'≠')^moxpr[parent]]
parent[redox]∩parent[idxpr_db]

```

```

:Namespace
  DB←LD'DYNA'
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]
  +/ (QUERY_MASK≠DB)[;AGES]
:EndNamespace

```

```

db←_defs^(nse roots)^named'DB'
dbex←(type=Var)^(vardef∈db)^expr[parent]
{dbex[parent]v←dbex^expr[parent]}*≡θ
idxpr←parent[_prmf n^(named'[')^dexpr[parent]]
idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]
redox←parent[_prmop^(named'≠')^moxpr[parent]]
parent[redox]∩parent[idxpr_db]

```

```

:Namespace
  DB←LD'DYNA'
  QUERY_MASK←DB[;EMAIL]∈(LD'DYALOG')[;EMAIL]
  +/(QUERY_MASK≠DB)[;AGES]
:EndNamespace

```

```

db←_defs^(nse roots)^named'DB'
dbex←(type=Var)^(vardef∈db)^expr[parent]
{dbex[parent]v←dbex^expr[parent]}*≡θ
idxpr←parent[_prmf n^(named'[')^dexpr[parent]]
idxpr_db←parent[_dbex^(≠parent)^parent∈idxpr]
redox←parent[_prmop^(named'≠')^moxpr[parent]]
parent[redox]∩parent[idxpr_db]

```

We could refactor...

- `(QUERY_MASK≠DATABASE)[;AGES]`
- `QUERY_MASK≠DATABASE[;AGES]`
- `AGES⇒QUERY_MASK◦≠DATABASE`
- C.f. – Asher on Trees/Parent Vectors

```
QUERY_SAVE ← { 'Q1' STORE  $\alpha \neq \omega$  }  
RESULTS ← MASK QUERY_SAVE DATABASE  
AVG_AGES ← ( +  $\neq$  ÷  $\neq$  ) ( LOAD 'Q1' ) [ ; AGES ]
```

```

QUERY_SAVE ← { 'Q1' STORE  $\alpha \neq \omega$  }
RESULTS ← MASK QUERY_SAVE DATABASE
AVG_AGES ← (+  $\neq \div$ ) (LOAD 'Q1') [ ; AGES ]

```

```

db ←  $\underline{1}$  (ns ∈ roots) ∧ defs ∧ named 'DB'
dbex ← (type = Var) ∧ (vardef ∈ db) ∧ expr[parent]
{ dbex[parent] v ← dbex ∧ expr[parent] } * ≡  $\theta$ 
idxpr ← parent[ $\underline{1}$ prmf n ∧ (named' [' ) ∧ dexpr[parent]]
idxpr_db ← parent[ $\underline{1}$ dbex ∧ ( $\neq$ parent) ∧ parent ∈ idxpr]
redox ← parent[ $\underline{1}$ pr mop ∧ (named'  $\neq$ ' ) ∧ moxpr[parent]]
parent[redox] ∩ parent[idxpr_db]

```

```

QUERY_SAVE ← { 'Q1' STORE  $\alpha \neq \omega$  }
RESULTS ← MASK QUERY_SAVE DATABASE
AVG_AGES ← (+  $\neq \div$ ) (LOAD 'Q1') [ ; AGES ]

```

```

db ←  $\perp$  (ns ∈ roots) ∧ defs ∧ named 'DB'
dbex ← (type = Var) ∧ (vardef ∈ db) ∧ expr[parent]
{ dbex[parent] v ← dbex ∧ expr[parent] } * ≡  $\theta$ 
idxpr ← parent[ $\perp$ prmf n ∧ (named' [' ) ∧ dexpr[parent]]
idxpr_db ← parent[ $\perp$ dbex ∧ ( $\neq$ parent) ∧ parent ∈ idxpr]
redox ← parent[ $\perp$ prmop ∧ (named'  $\neq$ ' ) ∧ moxpr[parent]]
parent[redox] ∩ parent[idxpr_db]

```

```
QUERY_SAVE ← { 'Q1' STORE  $\alpha \neq \omega$  }  
RESULTS ← MASK QUERY_SAVE DB  
AVG_AGES ← (+  $\neq \div$ ) (LOAD 'Q1') [ ; AGES ]
```

```
db ←  $\underline{1}$  (ns ∈ roots) ∧ defs ∧ named 'DB'  
dbex ← expr ∧ uses v . ∧ used ∈ db
```

```
idxpr ← parent [  $\underline{1}$  prmf n ∧ (named ' ( ' ) ∧ dexpr [parent] ]  
idxpr_db ← parent [  $\underline{1}$  dbex ∧ (≠ parent) ∧ parent ∈ idxpr ]  
redox ← parent [  $\underline{1}$  pr mop ∧ (named ' ≠ ' ) ∧ moxpr [parent] ]  
 $\underline{1}$  (uses ∈ redox) ∧ uses v . ∧ used ∈ idxpr_db
```

```
QUERY_SAVE ← { 'Q1' STORE  $\alpha \neq \omega$  }  
RESULTS ← MASK QUERY_SAVE DB  
AVG_AGES ← (+  $\neq \div$ ) (LOAD 'Q1') [ ; AGES ]
```

```
db ←  $\lambda$  (ns ∈ roots) ^ defs ^ named 'DB'  
dbex ← expr ^ uses v . ^ used ∈ db
```

```
idxpr ← parent [  $\lambda$  pr m fn ^ (named ' [ ' ) ^ dexpr [ parent ] ]  
idxpr_db ← parent [  $\lambda$  dbex ^ (  $\neq$  parent ) ^ parent ∈ idxpr ]  
redox ← parent [  $\lambda$  pr mop ^ (named '  $\neq$  ' ) ^ moxpr [ parent ] ]  
 $\lambda$  (uses ∈ redox) ^ uses v . ^ used ∈ idxpr_db
```

```

QUERY_SAVE ← { 'Q1' STORE  $\alpha \neq \omega$  }
RESULTS ← MASK QUERY_SAVE DB
AVG_AGES ← (+  $\neq \div$ ) (LOAD 'Q1') [ ; AGES ]

```

```

db ←  $\lambda$  (ns  $\in$  roots)  $\wedge$  defs  $\wedge$  named 'DB'
dbex ← expr  $\wedge$  uses v .  $\wedge$  used  $\in$  db

```

```

idxpr ← parent [ $\lambda$  pr mfn  $\wedge$  (named' [' )  $\wedge$  dexpr [parent]]
idxpr_db ← parent [ $\lambda$  dbex  $\wedge$  ( $\neq$  parent)  $\wedge$  parent  $\in$  idxpr]
redox ← parent [ $\lambda$  pr mop  $\wedge$  (named'  $\neq$  ' )  $\wedge$  moxpr [parent]]
 $\lambda$  (uses  $\in$  redox)  $\wedge$  uses v .  $\wedge$  used  $\in$  idxpr_db

```

```
QUERY_SAVE ← { 'Q1' STORE  $\alpha \neq \omega$  }  
RESULTS ← MASK QUERY_SAVE DB  
AVG_AGES ← ( +  $\neq \div$  ) ( LOAD 'Q1' ) [ ; AGES ]
```

```
db ← l ( ns ∈ roots ) ^ defs ^ named ' DB '  
dbex ← expr ^ uses v . ^ used ∈ db
```

```
idxpr ← parent [ l prmf n ^ ( named ' [ ' ) ^ dexpr [ parent ] ]  
idxpr_db ← parent [ l dbex ^ (  $\neq$  parent ) ^ parent ∈ idxpr ]  
redox ← parent [ l pr mop ^ ( named '  $\neq$  ' ) ^ moxpr [ parent ] ]  
l ( uses ∈ redox ) ^ uses v . ^ used ∈ idxpr_db
```

```
QUERY_SAVE ← { 'Q1' STORE  $\alpha \neq \omega$  }  
RESULTS ← MASK QUERY_SAVE DB  
AVG_AGES ← (+÷≠) (LOAD 'Q1') [ ; AGES ]
```

```
db ← ⊥ (ns ∈ roots) ∧ defs ∧ named 'DB'  
dbex ← expr ∧ uses v . ∧ used ∈ db
```

```
idxpr ← parent [ ⊥ pr mfn ∧ (named ' ( ' ) ∧ dexpr [ parent ] ]  
idxpr_db ← parent [ ⊥ dbex ∧ (≠ parent) ∧ parent ∈ idxpr ]  
redox ← parent [ ⊥ pr mop ∧ (named ' ≠ ' ) ∧ moxpr [ parent ] ]  
⊥ (uses ∈ redox) ∧ uses v . ∧ used ∈ idxpr_db
```

```
QUERY_SAVE ← { 'Q1' STORE  $\alpha \neq \omega$  }  
RESULTS ← MASK QUERY_SAVE DB  
AVG_AGES ← ( + ÷ ≠ ) ( LOAD 'Q1' ) [ ; AGES ]
```

```
db ← ⊥ ( ns ∈ roots ) ∧ defs ∧ named 'DB'  
dbex ← expr ∧ uses v . ∧ used ∈ db
```

```
idxpr ← parent [ ⊥ prmf n ∧ ( named ' [ ' ) ∧ dexpr [ parent ] ]  
idxpr_db ← parent [ ⊥ dbex ∧ ( ≠ parent ) ∧ parent ∈ idxpr ]  
redox ← parent [ ⊥ pr mop ∧ ( named ' ≠ ' ) ∧ moxpr [ parent ] ]  
⊥ ( uses ∈ redox ) ∧ uses v . ∧ used ∈ idxpr_db
```

Static Analysis Tool for APL Notation

Special Analyses

Code Quality Rules

Security Rules

Tracing data flow for you

Migration Assistance

A set of common rules and checks

Security Vulnerabilities

Support Our Customers around Compliance and Code Quality Audits

Standardization of Security Habits

Integration with Industry Products

(SonarQube, GitGuardian, Coverity, *etc.*)

Make things easier...

Easy customization

Easy application integration

Easy training, onboarding

Easy deployment

Provide your customers
with additional confidence and trust

Thank You.

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brandon@dyalog.com

<https://github.com/Co-dfns/Co-dfns>