

Language Features of version 18.0 in Depth

Adám Brudzewsky



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First in a series!

New

<code>□C</code>	Case convert
<code>fög</code>	Over
<code>fög</code>	Atop
<code>≠Y</code>	Unique mask
<code>A~</code>	Constant
<code>□DT</code>	Date-time
<code>1200I</code>	Format date-time

Improved

<code>□JSON@</code>	'HighRank'
<code>□JSON@</code>	'Dialect'
<code>□R/□S@</code>	'Regex'
<code>□INPUT@</code>	'NEOL'
<code>⌊Y</code>	
<code>X<Y</code>	
<code>↑[k]Y</code>	

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<code>↑[k]Y</code>

Case Convert

makes 819I obsolete

□C

Wait, what?

Uppercase: 1 (8 1 9 I) Y ⇒ 1 □ C Y

Lowercase: 8 1 9 I Y ⇒ □ C Y

Wait, what?

Uppercase: 1 (8 1 9 I) Y ⇒ 1 □ C Y

Lowercase: 8 1 9 I Y ⇒ □ C Y

Lowercase: 0 (8 1 9 I) Y ⇒ □ C Y

Wait, what?

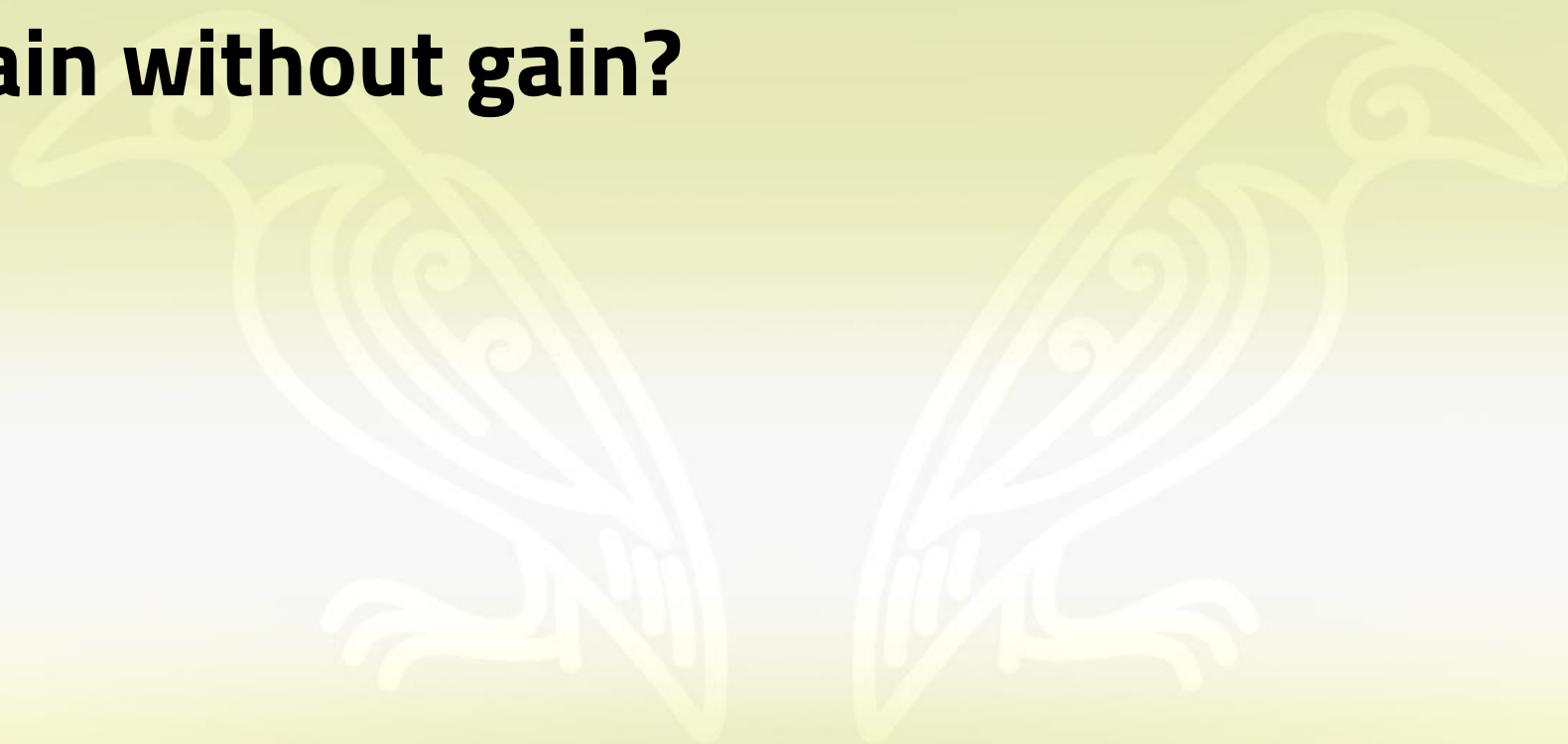
Uppercase: $1 (819I) Y \Rightarrow 1 \square C Y$

Lowercase: $819I Y \Rightarrow \square C Y$

Lowercase: $0 (819I) Y \Rightarrow \square C Y$

Big $\leftarrow \{ \alpha \leftarrow 0$
 $\alpha : 1 \square C \omega$
 $\square C \omega \}$

Pain without gain?



Pain without gain?

```
819I 'Hi' #(3J14 'PI')
```

DOMAIN ERROR: Invalid right argument

```
819I 'Hi' #(3J14 'PI')
```

^

Pain without gain?

```
819I'Hi'#(3J14'PI')
```

```
DOMAIN ERROR: Invalid right argument
```

```
819I'Hi'#(3J14'PI')
```

```
^
```

```
□C'Hi'#(3J14'PI')
```

```
hi # 3J14 pi
```

Pain without gain?

```
+ 'Hi' # (3J14 'PI')
```

```
Hi # 3J-14 PI
```

```
□ C 'Hi' # (3J14 'PI')
```

```
hello # 3J14 pi
```

Case Convert

Monadic `⊞C`: Case Fold
normalisation
for machine comparison

Dyadic `⊞C`: Case Map
display form
for human readers

Case Folding: ☐C Y

907

ABCDEFGHIJKLMNOPQRSTUVWXYZ

abcdefghijklmnopqrstuvwxyz

Case Folding: ☐C Y



Case Mapping: X C Y

1 : Upper

Origin

-1 : Lower

A	←	A a	⇒	a
B	←	B b	⇒	b
C	←	C c	⇒	c
D	←	D d	⇒	d
E	←	E e	⇒	e
F	←	F f	⇒	f
G	←	G g	⇒	g

Folding vs Mapping



'Μωσής'

Μωσής

□C 'Μωσής'

A fold

μωσής

⁻¹□C 'Μωσής'

A map

μωσής

Folding vs Mapping

\square_C 'Μωυσής' 'ΜΩΥΣΉΣ' A fold
 μωυσ^ήσ μωυσ^ήσ
 $^{-1}$ \square_C 'Μωυσής' 'ΜΩΥΣΉΣ' A map
 μωυσ^ής μωυσ^ήσ

Folding vs Mapping

	□C	'Μωυσής'	'ΜΩΥΣΉΣ'	Α fold	
μωυ	σή	σ	μωυ	σή	σ
1	□C	'Μωυσής'	'ΜΩΥΣΉΣ'	Α map	
ΜΩΥ	Σ	Ή	Σ	Ή	Σ

Folding vs Mapping

*"Street"
in German*

	□C	'Straße'	'STRASSE'	A fold
straße		straße		
1	□C	'Straße'	'STRASSE'	A map
STRASSE		STRASSE		

Folding vs Mapping

*"Street"
in German*


□C 'Straße' 'STRASSE' A fold
 strasse strasse

1 □C 'Straße' 'STRASSE' A map
 STRASSE STRASSE

STRASSE STRASSE

Would you ever map?

- All-lowercase to generate URLs or hash-tags
'Ο Μωσής Ζει' ⇒ '#ομωσήςζει'
- All-caps for display purposes
'Sale' ⇒ 'S A L E'
- Title-case a heading...
'Would you ever map?' ⇒
'Would You Ever Map?'



```
t ← 'Would you ever map?'
```

```
t←'Would you ever map?'  
-1*0,-1↓' ≠t
```

```
1 -1 -1 -1 -1 -1 1 -1 -1 -1 1 -1 -1 -1 -1 1 -1 -1 -1
```



```
t←'Would you ever map?'
-1*0,-1↓' ≠t
1 -1 -1 -1 -1 -1 1 -1 -1 -1 1 -1 -1 -1 -1 1 -1 -1 -1
(-1*0,-1↓' ≠t)⊞C t
```

```

t←'Would you ever map?'
-1*0,-1↓' ≠t
1 -1 -1 -1 -1 -1 1 -1 -1 -1 1 -1 -1 -1 -1 1 -1 -1 -1
(-1*0,-1↓' ≠t)□C t
DOMAIN ERROR: Invalid left argument
(-1*0,-1↓' ≠t)□C t
      ^

```

```

t←'Would you ever map?'
-1*0,-1↓' ≠t
1 -1 -1 -1 -1 -1 1 -1 -1 -1 1 -1 -1 -1 -1 1 -1 -1 -1
(-1*0,-1↓' ≠t)□C t
DOMAIN ERROR: Invalid left argument
(-1*0,-1↓' ≠t)□C t
      ^
(-1*0,-1↓' ≠t)□C t
Would You Ever Map?

```

New

$\square C$	Case convert
$f \ddot{o} g$	Over
$f \circ g$	Atop
$\neq Y$	Unique mask
$A \ddot{\sim}$	Constant
$\square DT$	Date-time
$1200 \mp$	Format date-time

Improved

$\square JSON \oplus 'HighRank'$
$\square JSON \oplus 'Dialect'$
$\square R / \square S \oplus 'Regex'$
$\square NPUT \oplus 'NEOL'$
$\underline{z} Y$
$X < Y$
$\uparrow [k] Y$

Compositions

We have a few compositions...

$$f \circ g \ \omega \quad \Leftrightarrow \quad f \quad g \ \omega$$

$$(f \ g) \ \omega \quad \Leftrightarrow \quad f \quad g \ \omega$$

$$\alpha \ f \circ g \ \omega \quad \Leftrightarrow \quad \alpha \ f \quad g \ \omega$$

$$\alpha (f \ g) \ \omega \quad \Leftrightarrow \quad f \ \alpha \ g \ \omega$$

We have a few compositions... but is there a system to these?

$$f \circ g \ \omega \quad \Leftrightarrow \quad f \quad g \ \omega$$

$$(f \ g) \ \omega \quad \Leftrightarrow \quad f \quad g \ \omega$$

$$\alpha \ f \circ g \ \omega \quad \Leftrightarrow \quad \alpha \ f \quad g \ \omega$$

$$\alpha (f \ g) \ \omega \quad \Leftrightarrow \quad f \ \alpha \ g \ \omega$$



f g

g f

g f g

fY gY	<pre>graph BT; g --> f; omega --> g;</pre>		
fY XgY	<pre>graph BT; g --> f; alpha --> g; omega --> g;</pre>		
XfY gY	<pre>graph BT; alpha --> f; g --> f; omega --> g;</pre>	<pre>graph BT; g --> f; alpha --> g; omega --> f;</pre>	<pre>graph BT; g --> f; alpha --> g; omega --> g; g --> f;</pre>



	f g	g f	g f g
fY gY	$f \circ g$ 		
fY XgY			
XfY gY	$f \circ g$ 	$f \overset{\sim}{\circ} g \overset{\sim}{\circ}$ 	



	f g	g f	g f g
fY gY	$f \ddot{g}$ $f \circ g$ $f \ddot{g}$ 		
fY XgY	$f \ddot{g}$ 		
XfY gY	$f \circ g$ 	$f \ddot{\circ} g \ddot{\circ}$ 	$f \ddot{g}$

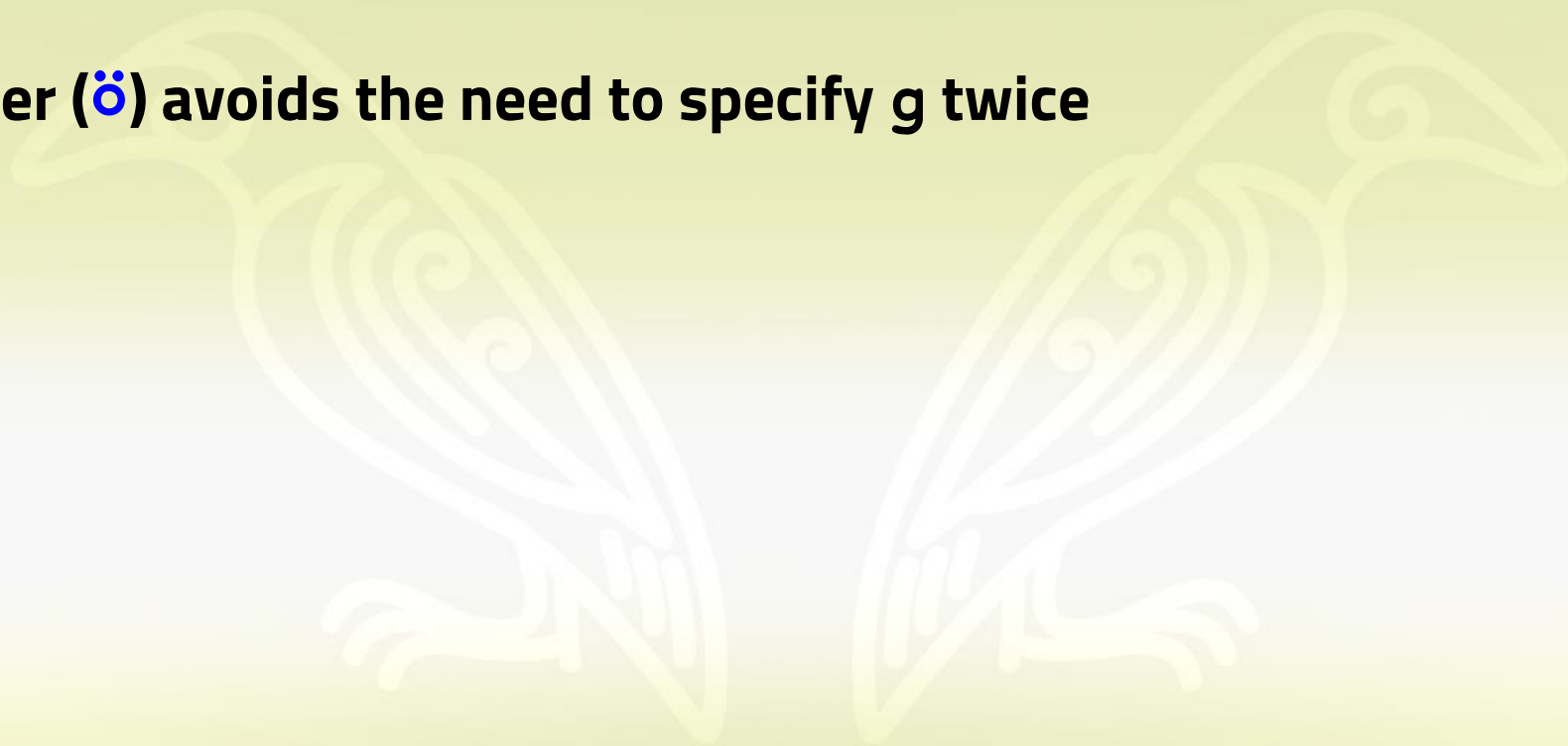


Over

on your keyboard as APL+Shift+Over (Shift+o)

ö

Over (ö) avoids the need to specify g twice



Over (ö) avoids the need to specify g twice

$$\{M \leftarrow + / \div \neq \diamond (M \alpha) + (M \omega)\}$$

Over (ö) avoids the need to specify g twice

$$\{M \leftarrow + / \div \neq \diamond (M \alpha) + (M \omega)\}$$

$$((+ / \div \neq) \dashv) + ((+ / \div \neq) \vdash)$$

Over (\ddot{o}) avoids the need to specify g twice

$$\{M \leftarrow + / \div \neq \diamond (M \alpha) + (M \omega)\}$$

$$((+ / \div \neq) \dashv) + ((+ / \div \neq) \vdash)$$

$$(M \dashv) + \circ (M \leftarrow + / \div \neq)$$

Over ($\ddot{\circ}$) avoids the need to specify g twice

$$\{M \leftarrow + / \div \neq \diamond (M \alpha) + (M \omega)\}$$

$$((+ / \div \neq) \rightarrow) + ((+ / \div \neq) \vdash)$$

$$(M \rightarrow) + \circ (M \leftarrow + / \div \neq)$$

$$+ \ddot{\circ} (+ / \div \neq)$$

Over (ö) represents a very common construct

Sum of squares

+ö(*o2)

Caseless function application

≡ö□C ιö□C εö□C etc.

Anagram?

≡ö{ω[Δω]~' '}

Juxtaposition ({α ω})

,öc

Laminate vectors (, [1.5])

,ö;

Over (ö) is closely related to Beside (•)

$$f \bullet g \ \omega \iff f \ (g \ \omega)$$

$$f \ddot{o} g \ \omega \iff f \ (g \ \omega)$$

$$\alpha \ f \bullet g \ \omega \iff \alpha \ f \ (g \ \omega)$$

$$\alpha \ f \ddot{o} g \ \omega \iff (g \ \alpha) \ f \ (g \ \omega)$$

Over (ö) is closely related to Beside (•)



*New
name!*

$$f \bullet g \ \omega \iff f \ (g \ \omega)$$

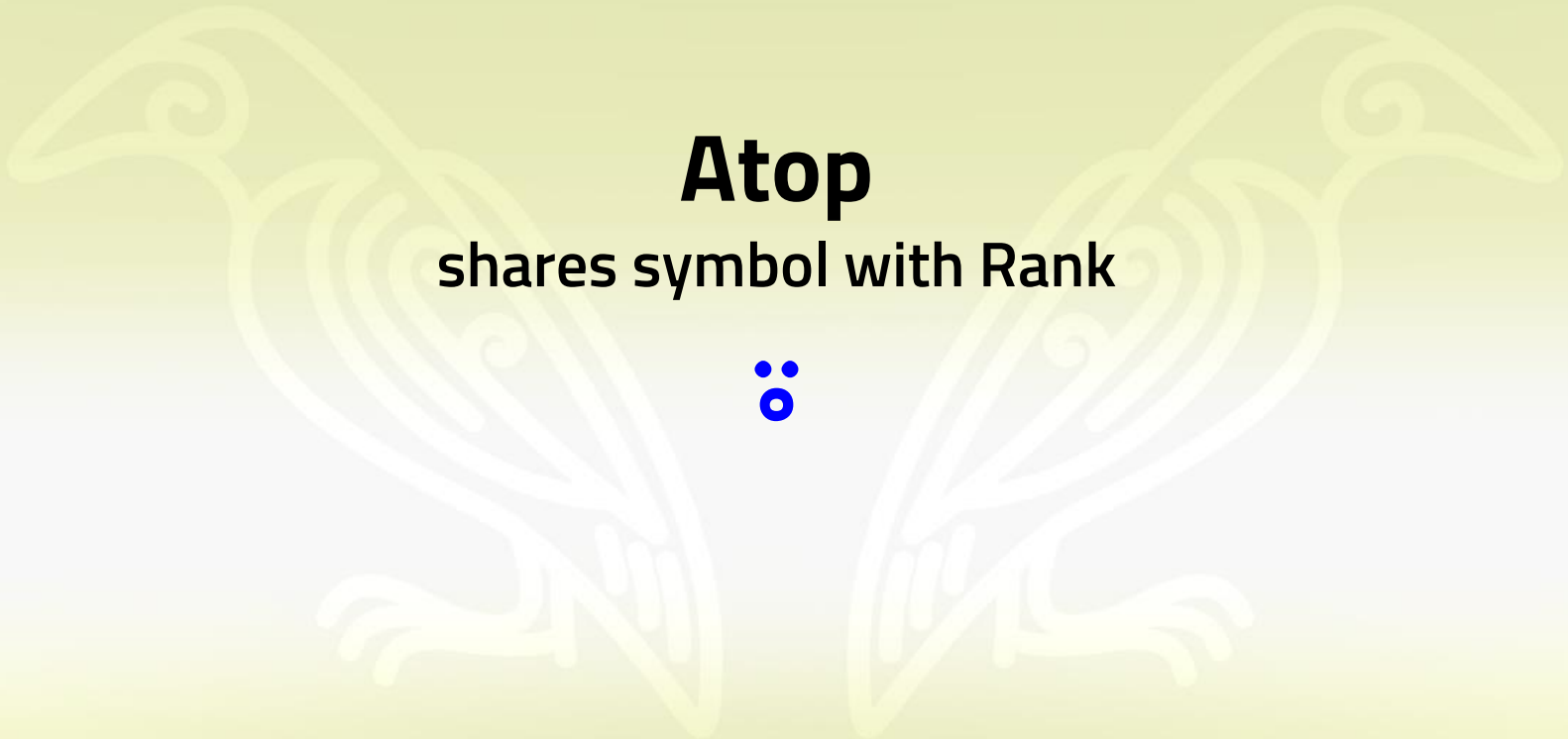
$$f \ddot{o} g \ \omega \iff f \ (g \ \omega)$$

$$\alpha \ f \bullet g \ \omega \iff \alpha \ f \ (g \ \omega)$$

$$\alpha \ f \ddot{o} g \ \omega \iff (g \ \alpha) \ f \ (g \ \omega)$$

Enough, give it to me already!

Over $\leftarrow \{ (\omega\omega \ \alpha) \ \alpha\alpha \ (\omega\omega \ \omega) \}$



Atop
shares symbol with Rank

ö

Atop ($\ddot{\circ}$) is closely related to Beside (\circ)

$f \circ g \ \omega \iff f \quad g \ \omega$

$f \ddot{\circ} g \ \omega \iff f \quad g \ \omega$

$\alpha \ f \circ g \ \omega \iff \alpha \ f \quad g \ \omega$

$\alpha \ f \ddot{\circ} g \ \omega \iff f \quad \alpha \ g \ \omega$

Wait a minute! — We already have atop (())

$(f\ g)\omega \iff f\ g\ \omega$

$f\ddot{\circ}g\ \omega \iff f\ g\ \omega$

$\alpha(f\ g)\omega \iff \textcircled{f\ \alpha}\ g\ \omega$

$\alpha\ f\ddot{\circ}g\ \omega \iff \textcircled{f\ \alpha}\ g\ \omega$

- mitigates function/operator duality by forcing function-ness

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$$\{ \ / \ / \ \square \vee \text{FI} \ \omega \} \Leftrightarrow \{ (\ / \ /) \ \square \vee \text{FI} \ \omega \}$$

• mitigates function/operator duality by forcing function-ness

$$\begin{array}{l}
 \{ \quad / \quad / \quad \square \forall F I \quad \omega \} \\
 \{ \supset \quad / \quad / \quad \square \forall F I \quad \omega \}
 \end{array}
 \Leftrightarrow
 \begin{array}{l}
 \{ \quad (/ \quad /) \quad \square \forall F I \quad \omega \} \\
 \{ (\supset \quad /) \quad / \quad \square \forall F I \quad \omega \}
 \end{array}$$

• mitigates function/operator duality by forcing function-ness

$$\begin{array}{l}
 \{ \quad / \quad / \quad \Box VFI \quad \omega \} \quad \Leftrightarrow \quad \{ \quad (/ \quad /) \quad \Box VFI \quad \omega \} \\
 \{ \supset \quad / \quad / \quad \Box VFI \quad \omega \} \quad \Leftrightarrow \quad \{ (\supset \quad /) \quad / \quad \Box VFI \quad \omega \} \\
 \{ \supset \quad \vdash \ddot{\circ} / \quad / \quad \Box VFI \quad \omega \} \quad \Leftrightarrow \quad \{ \supset (/ \quad /) \quad \Box VFI \quad \omega \}
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 \{ \supset \quad \vdash \ddot{\circ} / \quad / \quad \Box \forall F I \quad \omega \} \Leftrightarrow \{ \supset (/ \quad /) \quad \Box \forall F I \quad \omega \}
 \end{array}$$

$$\{ (5 < \omega) \quad / \quad \omega \} \Leftrightarrow \{ (5 < \omega) \quad / \quad \omega \}$$

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 \{ \quad / \quad / \quad \Box VFI \quad \omega \} \Leftrightarrow \{ \quad (/ \quad /) \quad \Box VFI \quad \omega \} \\
 \{ \supset \quad / \quad / \quad \Box VFI \quad \omega \} \Leftrightarrow \{ (\supset \quad /) \quad / \quad \Box VFI \quad \omega \} \\
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 \end{array}$$

$$\begin{array}{l}
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 ((5 < \vdash) \quad / \quad \vdash) \Leftrightarrow \{ (\{ 5 < \omega \} \quad /) \quad \omega \}
 \end{array}$$

• mitigates function/operator duality by forcing function-ness

$$\begin{array}{l}
 \{ \quad / \quad / \quad \Box VFI \quad \omega \} \Leftrightarrow \{ \quad (/ \quad /) \quad \Box VFI \quad \omega \} \\
 \{ \supset \quad / \quad / \quad \Box VFI \quad \omega \} \Leftrightarrow \{ (\supset \quad /) \quad / \quad \Box VFI \quad \omega \} \\
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 \end{array}$$

$$\begin{array}{l}
 \{ (5 < \omega) \quad / \quad \omega \} \Leftrightarrow \{ (5 < \omega) \quad / \quad \omega \} \\
 ((5 < \vdash) \quad / \quad \vdash) \Leftrightarrow \{ (\{ 5 < \omega \} \quad /) \quad \omega \} \\
 ((5 < \vdash) \quad \vdash \ddot{\circ} / \quad \vdash) \Leftrightarrow \{ (5 < \omega) \quad / \quad \omega \}
 \end{array}$$

... and make the structure of tacit functions clearer

```
{(Pea Que α Are ω)Ess(Tee You α Vie ω)}
```

... and make the structure of tacit functions clearer

```
{(Pea Que α Are ω)Ess(Tee You α Vie ω)}  
((Pea Que)Are)Ess((Tee You)Vie)
```

... and make the structure of tacit functions clearer

```
{(Pea Que α Are ω)Ess(Tee You α Vie ω)}
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```
((Pea Que)Are)Ess((Tee You)Vie)
```

```
(Pea◦Que Are)Ess(Tee◦You Vie)
```


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```
(Pea◦Que Are)Ess◦Tee◦You Vie
```

... and make the structure of tacit functions clearer

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((Pea Que)Are)Ess((Tee You)Vie)
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(Pea◦Que Are)Ess(Tee◦You Vie)
```

```
(Pea◦Que Are)Ess◦Tee◦You Vie
```

```
Pea◦◦Que◦◦Are Ess Tee◦◦You◦◦Vie
```

Enough, give it to me already!

Atop $\leftarrow \{ \alpha \leftarrow \vdash \quad \diamond \quad \alpha\alpha \quad \alpha \quad \omega\omega \quad \omega \}$

Unique Mask

(a.k.a. nub-sieve)

≠ Y

u 'Mississippi'
Misp



```
u 'Mississippi'  
Misp  
≠ 'Mississippi'  
1 1 1 0 0 0 0 0 1 0 0
```

```
u 'Mississippi'
Misp
≠ 'Mississippi'
1 1 1 0 0 0 0 0 1 0 0
'Mississippi' ↑ö,öc ≠ 'Mississippi'
```

```

u 'Mississippi'
Misp
≠ 'Mississippi'
1 1 1 0 0 0 0 0 1 0 0
'Mississippi'   ↑ö, öc   ≠ 'Mississippi'
M i s s i s s i p p i
1 1 1 0 0 0 0 0 1 0 0

```



```

    u 'Mississippi'
Misp
    ≠ 'Mississippi'
1 1 1 0 0 0 0 0 1 0 0
    'Mississippi' ↑ö, öc ≠ 'Mississippi'
M i s s i s s i p p i
1 1 1 0 0 0 0 0 1 0 0
    (≠ ↑ö ↓) 'Mississippi'
Misp

```

Why, though?

$\neq Y$

is to

$\cup Y$

as

$\uparrow Y$

is to

Sort Y

▲ Y vs Sort Y

Sort 3 1 4 1 5

1 1 3 4 5

▲ 3 1 4 1 5

2 4 1 3 5

'Moses' [2 4 1 3 5]

oeMss

≠Y vs UY

U 3 1 4 1 5

3 1 4 5

≠ 3 1 4 1 5

1 1 1 0 1

1 1 1 0 1 / 'Moses'

Moss

Non-unique

u 'Mississippi'

Misp

≠ 'Mississippi'

1 1 1 0 0 0 0 0 1 0 0

~≠ 'Mississippi'

0 0 0 1 1 1 1 1 0 1 1

(~ö≠ **┌ö/┐** ┌) 'Mississippi'

sissippi

Enough, give it to me already!

Umask \leftarrow \sim = \neq

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- fög Over
- fög Atop
- ≠Y Unique mask
- A~ Constant
- DT Date-time
- 1200± Format date-time

Improved

- JSON: 'HighRank'
- JSON: 'Dialect'
- R/□S: 'Regex'
- INPUT: 'NEOL'
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Questions?

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<code>↑[k]Y</code>	

APL and Microsoft Excel

Richard Park



May 27, 2020